

# The Game of Life on Penrose Tilings: Pentagon Boat-Star

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## 1 Introduction

In the 1970s Roger Penrose and John Conway presented the mathematical world with two beautiful and intriguing developments - Penrose with his sets of tiles that tile the plane non-periodically providing us a class of aperiodic tilings known as Penrose tilings [6], and Conway with his *Game of Life*, a cellular automaton simulated on a lattice of square cells with extremely simple rules governing the future state of the cells that can lead to extremely complex “organisms” being born out of relatively simple starting configurations[5].

We have investigated what the result would be of playing the Game of Life on a Penrose tiling, and present findings from our simulations.

### 1.1 Definitions

We begin by defining some of the terms we shall be using throughout this document.

**Definition 1.** A *tiling of the plane*  $\mathfrak{T}$  is a countable union of closed sets (known as *tiles*)  $\mathfrak{T} = \{T_1, T_2, T_3, \dots\}$  which cover the entire plane  $\mathbb{R}^2$  without gaps or overlaps. In other words,  $\mathbb{R}^2 = \bigcup_{T \in \mathfrak{T}} T$ . [7][2]

**Definition 2.** A *non-periodic tiling* is a tiling without any translational symmetries.

**Definition 3.** A *patch* is a finite set of tiles in a tiling with the property that their union is a topological disk – in other words, is connected and simply connected, and cannot be disconnected by the deletion of a single point. [7]

## 2 The Game of Life

Conway’s Game of Life is traditionally played on an infinite grid of squares, with each square being known as a *cell*. The fate of a cell in the next generation of the simulation depends upon the states of its immediate neighbors as follows:[5]

1. If a cell is *alive* and has 2 or 3 live neighbors, it survives to the next generation. Otherwise, it dies.

2. If a cell is *dead* and has exactly 3 live neighbors, it becomes alive in the next generation.

These rules albeit simple, when applied to multiple certain starting patterns can result in the formation of remarkably complex figures. However, these figures or “organisms” as they are also called, can be mostly classified into a few broad categories. [1][5]

1. A *still life* is an organism in which all cells survive to the next generation and is thus in a state of equilibrium - no births, no deaths.
2. An *oscillator* is an organism that starts at some initial state and after a certain number of generations (known as its period) returns to that state in the same position.
3. A *spaceship* is an oscillator that translates across the plane by a certain vector in a specific number of generations.

Some examples of these patterns are presented below.

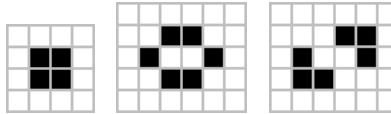


Figure 1: Examples of Still Lifes. (left to right) Block, Beehive and Aircraft carrier. Images from [8]



Figure 2: An oscillator, the blinker. Images from [8]

### 3 Penrose Tiling

Penrose tilings are created out of many different protosets with varying numbers of prototiles. However, here we will focus only on the P1 protoset, which consists of six prototiles as described below.

We use the substitution or inflate-and-subdivide method to generate the tiling. This involves taking a tile, enlarging it by a certain factor (which in our case happens to be the golden ratio  $\phi = \frac{1+\sqrt{5}}{2}$ ), and replace this enlarged tile by a certain arrangement of tiles with the same size as the original tile, and which collectively roughly imitate its shape. Applying this process repetitively  $n$  times we are able to produce an arbitrarily large patch which we call the  $n$ th level substitution patch.

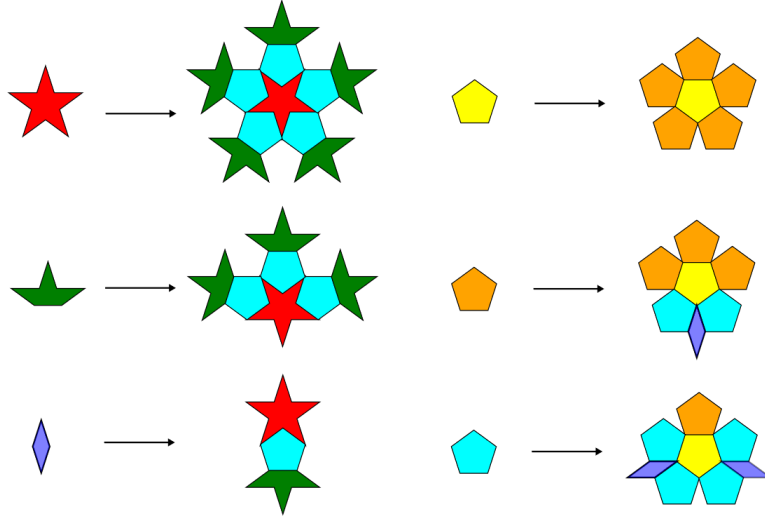


Figure 3: Substitution rules for the P1 Penrose prototile set. [7]

## 4 Playing the Game of Life on the Penrose Tiling

### 4.1 Generating Patches using Substitution Rules

We generate the tiling and data regarding the tiles and their neighbors which we then use to apply the Game of Life rules on the tiles. We choose to use the Julia programming language [3] for this purpose due to its easy learning curve and performance benefits, which helps greatly due to the large numbers of tiles we would be taking into consideration. We begin by defining a `Tile` struct which contains basic properties of a tile such as its type, its “origin” which is the first vertex of the tile to be drawn and from which all other vertices are calculated, the angle between its line of symmetry passing through the origin and the horizontal axis, and a dictionary for storing its neighbors. We observe that a tile can be formed using just 3 pieces of information - the type of tile, its origin and its angle with the horizontal. Let us take for example the case of a rhombus. Let us suppose we wish to draw a rhombus whose first vertex is at the origin of the plane and makes an angle of 0 with the horizontal. Then going anti-clockwise, we know that the vectors from the tile’s origin to the other points make angles of  $\frac{\pi}{10}$ , 0 and  $-\frac{\pi}{10}$ , and by calculating the magnitudes of these vectors, we can determine the coordinates for the other vertices. These coordinates can thus be generated on-the-fly later when we plot these tiles for a visual representation of them. Following this we define functions for the substitution of each type of tile (`substitute_star`, `substitute_boat`, `substitute_rhomb` etc.). We have determined the properties of the new tiles by calculations and some trial and error. Each function returns a tuple of new tiles.

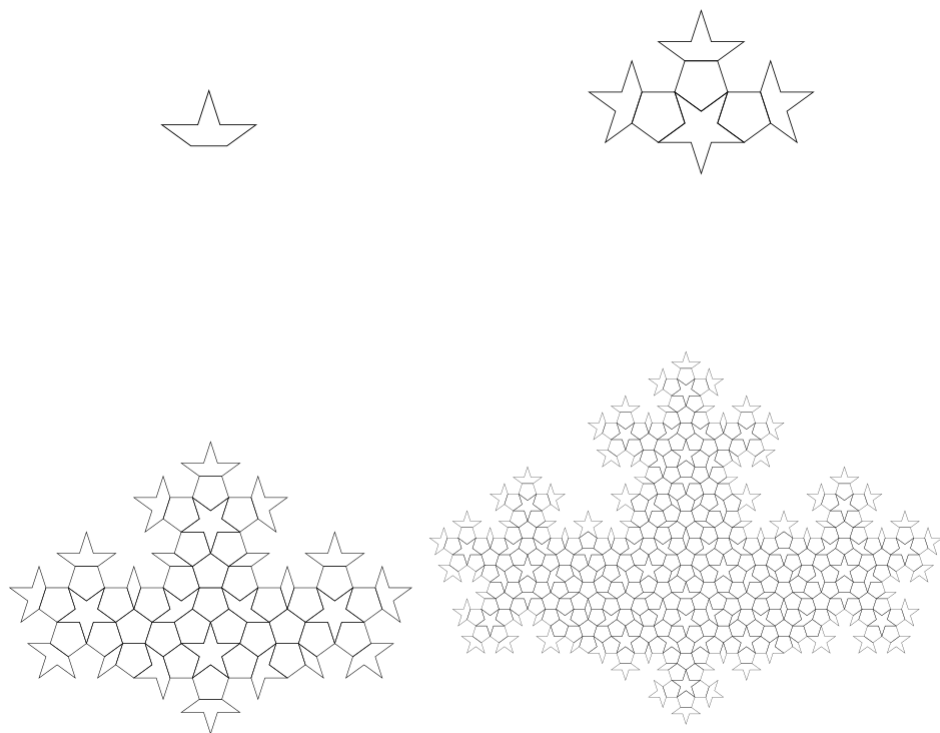


Figure 4: A few substitutions of the boat tile

We further define a more generalised function that reads the type of tile passed as an argument and returns the tiles yielded from the respective substitution function. We can then define a function `substitute_tiles` that takes as input an array of `Tile` objects and returns the array formed by aggregating the results of each of the tile’s substitutions. Applying this function  $n$  times to a given array of tiles will give us the  $n$ th level substitution patch from those tiles.

## 4.2 Game of Life Implementation

We begin by defining methods to determine the neighbors of a given tile in a patch. Choose a tile  $T$  and compute the coordinates of its vertices. Then, pick another tile  $T'$  and compute its vertices’ coordinates as well. We compare every vertex of  $T$  in counter-clockwise order to the vertices of  $T'$ . If any vertex of  $T$  is also a vertex of  $T'$ , then  $T'$  is a neighbor of  $T$ . If two such vertices exist, that must mean  $T$  and  $T'$  share an edge, and we say that  $T'$  is an *edge neighbor* of  $T$ . Otherwise,  $T'$  is said to be a *corner neighbor* of  $T$ . We set a numbering system for classifying the neighbors in the neighbor dictionary of the  $T$ . In this system we start numbering at 1, and the first  $s$  indices represent the edge neighbors corresponding to edges going anti-clockwise around the tile from the tile’s origin, where  $s$  is equal to the number of edges the tile has. Followed by this, the indices  $\{(s+t)|t \in \{1, 2, \dots, s\}\}$  hold the corner neighbors corresponding to the  $(s+t) - s = t$ th vertex. This process can very clearly take a lot of computation time due to the large amount of comparisons between points taking place. Thus, we compile the neighbor information of all tiles in a given patch into a dictionary whose keys are the indices of the tiles within the patch, and the values are their respective neighbor dictionaries. We store this dictionary into a JLD2 file which is a file format designed for saving and loading Julia data structures. This file can then be loaded whenever needed and the neighbor data used assuming that the underlying patch hasn’t been modified.

Now, since we have the number of neighbors for each tile in the patch, we implement the Game of Life rules for the same. First, we pass a list of live indices to initialize our simulation. Next, we extract the lists of neighbors for each tile and check if any of those neighbors’ indices are alive. Then, according to the number of live neighbors the tile is determined to have, we decide whether or not to add the tile’s index to the list of live indices for the next generation which is returned by our function. We can then repeat this process for as many generations as we like, passing in the output of the previous generation as an input for the next one.

## 4.3 Plotting and Visualization

To visualize the tiling and the Game of Life being played on it, we utilized the `Luxor` package. After getting the tiles in a patch we proceed to compute the vertices for each of them. We then use `Luxor`’s `poly` method to join those vertices to form a closed polygon. When playing the Game of Life, we add a black, filled polygon in the same manner as for the tiling at the coordinates

corresponding to the tile at the live index. We generate an SVG image for every generation using Luxor's `@svg` macro and animate and store it to a GIF file using the `@animate` macro.

## 5 Results

We limited our testing space to six patches pulled from the 5th substitution patch of a type-1 pentagon tile, centered roughly around the center of the patch. We did so in an effort to get a hold of a reasonably large enough patch for our simulations while minimizing the effects of considering tiles on the boundary of the patch, which lead to inaccurate results due to having an inaccurate number of live neighbors counted.

After running the Game of Life simulation multiple times on these patches we have observed numerous still lifes and oscillators. No spaceships have yet been discovered. The smallest still lifes and oscillators found comprised of 3 tiles each. Oscillator periods vary between 2 and 6. We have presented some of the organisms we have identified below.

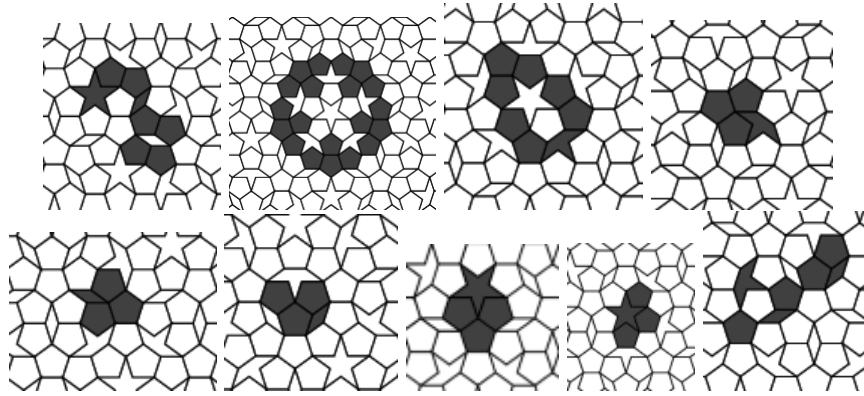


Figure 5: Still Lifes

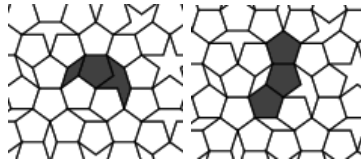


Figure 6: A period 2 oscillator

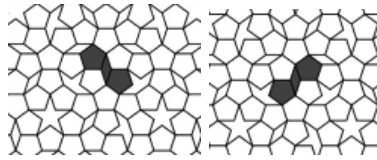


Figure 7: Another period 2 oscillator.

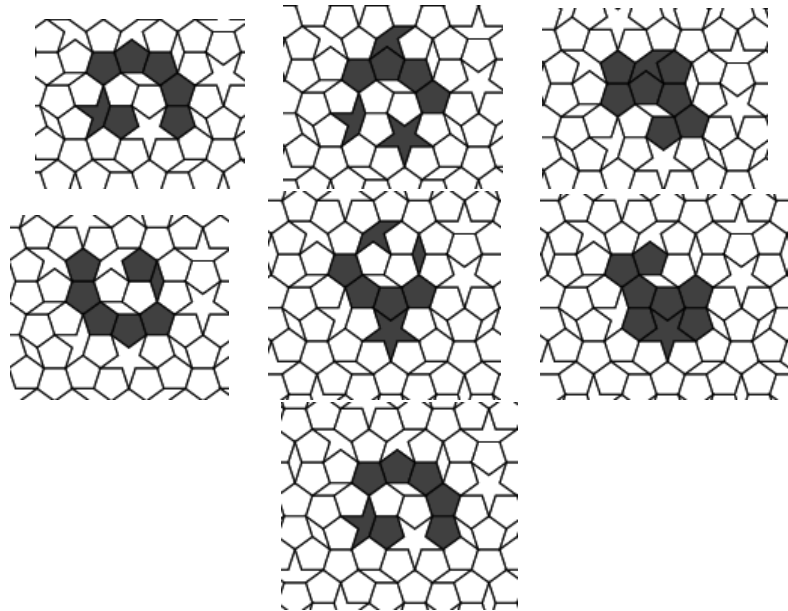


Figure 8: A period 6 oscillator

## References

- [1] *Lifewiki*, <https://conwaylife.com/wiki>.
- [2] C. ADAMS, *The tiling book*.
- [3] J. BEZANSON, A. EDELMAN, S. KARPINSKI, AND V. B. SHAH, *Julia: A fresh approach to numerical computing*, SIAM review, 59 (2017), pp. 65–98.
- [4] D. FRETTLÖH AND F. GÄHLER, *Penrose Pentagon Boat Star*, *Tiling Encyclopedia*, <https://tilings.math.uni-bielefeld.de/substitution/penrose-pentagon-boat-star/>.
- [5] M. GARDNER, *Mathematical games*, Scientific American, 223 (1970), p. 120–123.
- [6] ———, *Mathematical games*, Scientific American, 236 (1977), p. 110–121.
- [7] B. GRÜNBAUM AND G. C. SHEPHARD, *Tilings and Patterns, Second Edition*, Dover Publications, New York, 1987 (Reprint 2016).
- [8] N. JOHNSTON AND D. GREENE, *Conway's Game of Life: Mathematics and Construction*, Self-published, 2022, ch. Still Lives.