

# A Brief L<sup>A</sup>T<sub>E</sub>X Example\*

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MAX-HEAPIFY is an important subroutine for manipulating max-heaps. Its inputs are an array  $A$  and an index  $i$  into the array. When MAX-HEAPIFY is called, it is assumed that the binary trees rooted at  $\text{LEFT}(i)$  and  $\text{RIGHT}(i)$  are max-heaps, but that  $A[i]$  may be smaller than its children, thus violating the max-heap property. The function of MAX-HEAPIFY is to let the value at  $A[i]$  “float down” in the max-heap so that the subtree rooted at index  $i$  becomes a max-heap.

```
MAX-HEAPIFY( $A, i$ )
   $l \leftarrow \text{LEFT}(i)$ 
   $r \leftarrow \text{RIGHT}(i)$ 
  if  $l \leq \text{heap-size}[A]$  and  $A[l] > A[i]$ 
    then  $\text{largest} \leftarrow l$ 
    else  $\text{largest} \leftarrow i$ 
  if  $r \leq \text{heap-size}[A]$  and  $A[r] > A[\text{largest}]$ 
    then  $\text{largest} \leftarrow r$ 
  if  $\text{largest} \neq i$ 
    then exchange  $A[i] \leftrightarrow A[\text{largest}]$ 
    MAX-HEAPIFY( $A, \text{largest}$ )
```

Figure 1 illustrates the action of MAX-HEAPIFY. At each step, the largest of the elements  $A[i]$ ,  $A[\text{LEFT}(i)]$ , and  $A[\text{RIGHT}(i)]$  is determined, and its index is stored in  $\text{largest}$ . ...

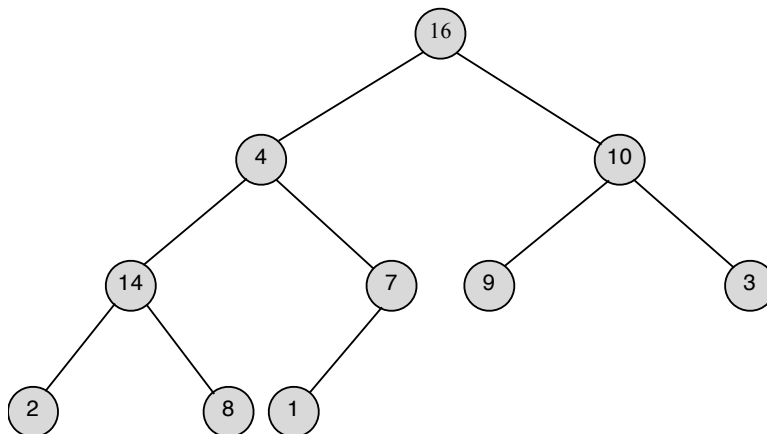


Figure 1: The action of MAX-HEAPIFY( $A, 2$ ), where  $\text{heap-size}[A] = 10$ . ...

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\*Taken from pp. 130–131 of *Introduction to Algorithms* by Cormen, et al.