Link-State Routing

Reading: Sections 4.2 and 4.3.4

CS 375: Computer Networks

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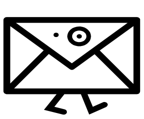
Goals of Today's Lecture

- Inside a router
  - –Control plane: routing protocols
  - –Data plane: packet forwarding
- Path selection
  - –Minimum-hop and shortest-path routing
  - –Dijkstra's algorithm
- Topology change
  - –Using beacons to detect topology changes
  - –Propagating topology information
- Routing protocol: Open Shortest Path First

What is Routing?

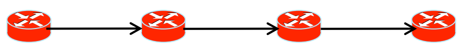
- A famous quotation from RFC 791

  "A *name* indicates what we seek.
  An *address* indicates where it is.
  A *route* indicates how we get there."
  -- Jon Postel

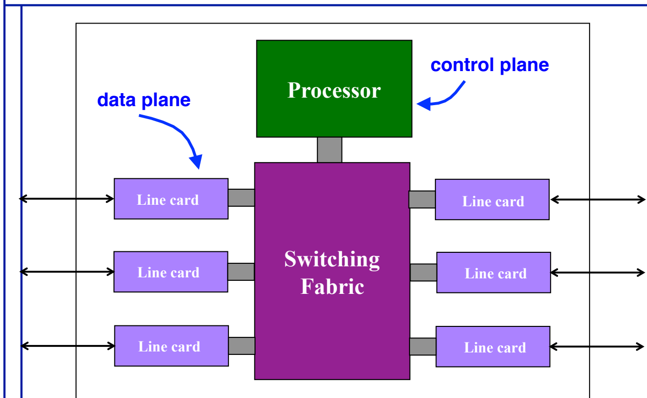


Routing vs. Forwarding

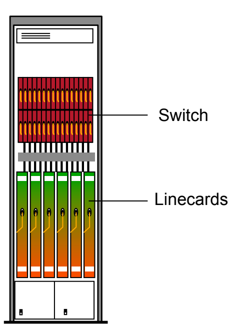
- Routing: control plane
  - –Computing paths the packets will follow
  - –Routers talking amongst themselves
  - –Individual router *creating* a forwarding table
- Forwarding: data plane
  - –Directing a data packet to an outgoing link
  - –Individual router *using* a forwarding table



Data and Control Planes



Router Physical Layout

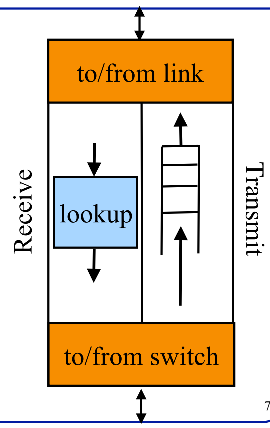


Juniper T series

Cisco 12000

Line Cards (Interface Cards, Adaptors)

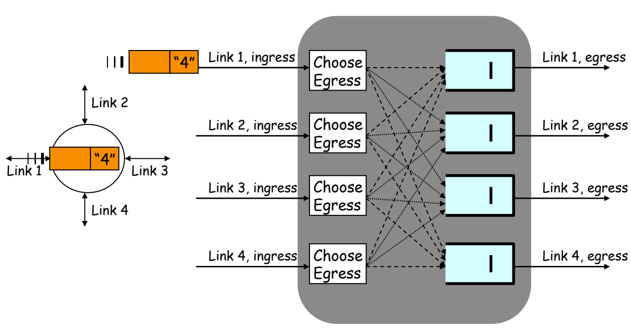
- Interfacing
  - Physical link
  - Switching fabric
- Packet handling
  - Packet forwarding
  - Decrement time-to-live
  - Buffer management
  - Link scheduling
  - Packet filtering
  - Rate limiting
  - Packet marking
  - Measurement



Switching Fabric

- Deliver packet inside the router
  - From incoming interface to outgoing interface
  - A small network in and of itself
- Must operate very quickly
  - Multiple packets going to same outgoing interface
  - Switch scheduling to match inputs to outputs
- Implementation techniques
  - Bus, crossbar, interconnection network, …
  - Running at a faster speed (e.g., 2X) than links
  - Dividing variable-length packets into fixed-size cells

Packet Switching



Router Processor

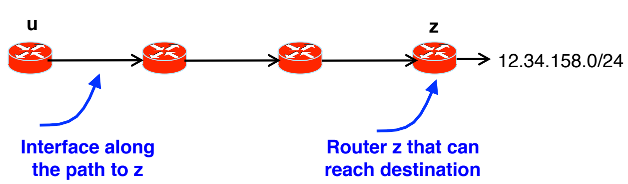
- So-called “Loopback” interface
  - IP address of the CPU on the router
- Interface to network administrators
  - Command-line interface for configuration
  - Transmission of measurement statistics
- Handling of special data packets
  - Packets with IP options enabled
  - Packets with expired Time-To-Live field
- Control-plane software
  - Implementation of the routing protocols
  - Creation of forwarding table for the line cards

Where do Forwarding Tables Come From?

- Routers have forwarding tables
  - Map IP prefix to outgoing link(s)
- Entries can be statically configured
  - E.g., “map 12.34.158.0/24 to Serial0/0.1”
- But, this doesn’t adapt
  - To failures
  - To new equipment
  - To the need to balance load
- That is where routing protocols come in

Computing Paths Between Routers

- Routers need to know two things
  - Which router to use to reach a destination prefix
  - Which outgoing interface to use to reach that router



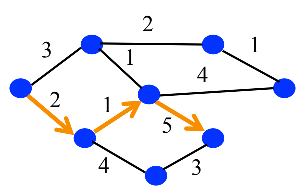
- Today’s class: just how routers reach each other
  - How u knows how to forward packets toward z

Computing the Shortest Paths

(assuming you already know the topology)

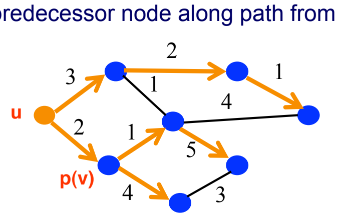
Shortest-Path Routing

- Path-selection model
  - Destination-based
  - Load-insensitive (e.g., static link weights)
  - Minimum hop count or sum of link weights



Shortest-Path Problem

- Given: network topology with link costs
  - **c(x,y)**: link cost from node x to node y
  - Infinity if x and y are not direct neighbors
- Compute: least-cost paths to all nodes
  - From a given source u to all other nodes
  - **p(v)**: predecessor node along path from source to v



Dijkstra's Shortest-Path Algorithm

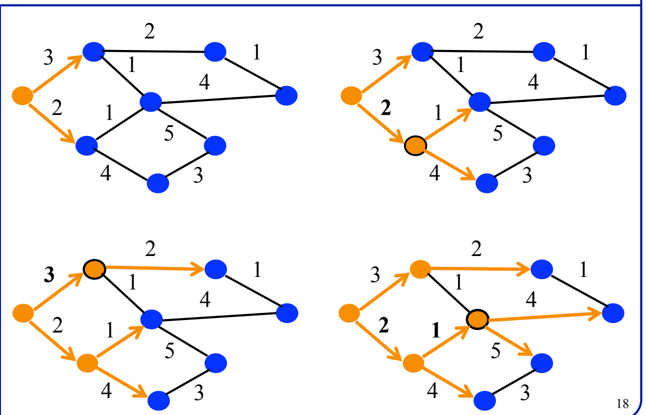
- Iterative algorithm
  - After k iterations, know least-cost path to k nodes
- **S**: nodes whose least-cost path definitively known
  - Initially, **S = {u}** where u is the source node
  - Add one node to S in each iteration
- **D(v)**: current cost of path from source to node v
  - Initially, **D(v) = c(u,v)** for all nodes v adjacent to u
  - … and **D(v) = ∞** for all other nodes v
  - Continually update D(v) as shorter paths are learned

Dijsktra's Algorithm

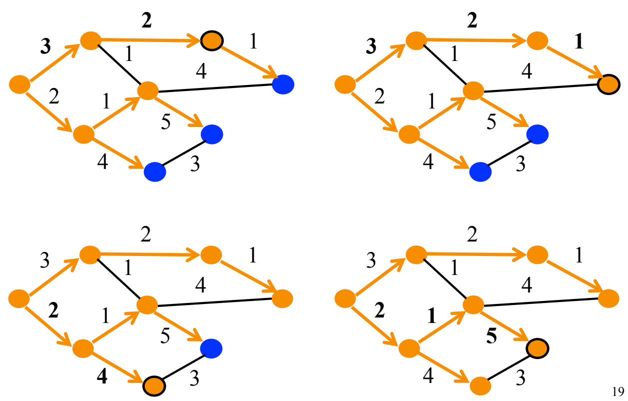
```
Initialization:
  S = {u}
  for all nodes v
    if (v is adjacent to u)
        D(v) = c(u,v)
    else D(v) = ∞

Loop
   find w not in S with the smallest D(w)
   add w to S
   update D(v) for all v adjacent to w and not in S:
       D(v) = min{D(v), D(w) + c(w,v)}
until all nodes in S
```


Dijkstra's Algorithm Example

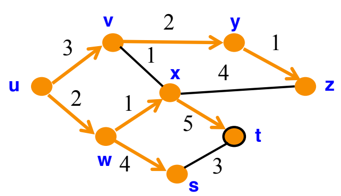


Dijkstra's Algorithm Example



Shortest-Path Tree

- Shortest-path tree from u
- Forwarding table at u



	link
v	(u,v)
w	(u,w)
x	(u,w)
y	(u,v)
z	(u,v)
s	(u,w)
t	(u,w)

Learning the Topology

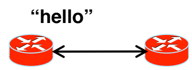
(by the routers talk amongst themselves)

Link-State Routing

- Each router keeps track of its incident links
  - Whether the link is up or down
  - The cost on the link
- Each router broadcasts the link state
  - To give every router a complete view of the graph
- Each router runs Dijkstra's algorithm
  - To compute the shortest paths
  - … and construct the forwarding table
- Example protocols
  - Open Shortest Path First (OSPF)
  - Intermediate System – Intermediate System (IS-IS)

Detecting Topology Changes

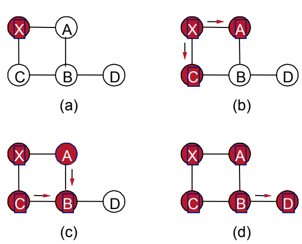
- Beaconing
  - Periodic "hello" messages in both directions
  - Detect a failure after a few missed "hellos"



- Performance trade-offs
  - Detection speed
  - Overhead on link bandwidth and CPU
  - Likelihood of false detection

Broadcasting the Link State

- Flooding
  - Node sends link-state information out its links
  - And then the next node sends out all of its links
  - … except the one where the information arrived



Broadcasting the Link State

- Reliable flooding
  - Ensure all nodes receive link-state information
  - … and that they use the latest version
- Challenges
  - Packet loss
  - Out-of-order arrival
- Solutions
  - Acknowledgments and retransmissions
  - Sequence numbers
  - Time-to-live for each packet

When to Initiate Flooding

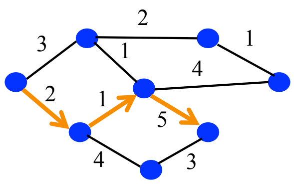
- Topology change
  - Link or node failure
  - Link or node recovery
- Configuration change
  - Link cost change
- Periodically
  - Refresh the link-state information
  - Typically (say) 30 minutes
  - Corrects for possible corruption of the data

When the Routers Disagree

(during transient periods)

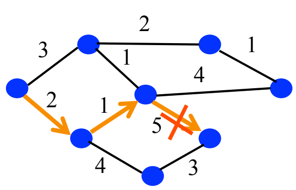
Convergence

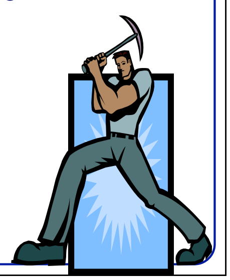
- Getting consistent routing information to all nodes
  - E.g., all nodes having the same link-state database
- Consistent forwarding after convergence
  - All nodes have the same link-state database
  - All nodes forward packets on shortest paths
  - The next router on the path forwards to the next hop



Transient Disruptions

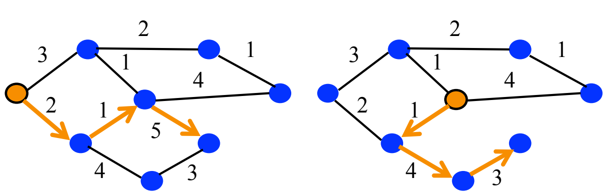
- Detection delay
  - A node does not detect a failed link immediately
  - … and forwards data packets into a “blackhole”
  - Depends on timeout for detecting lost hellos





Transient Disruptions

- Inconsistent link-state database
  - Some routers know about failure before others
  - The shortest paths are no longer consistent
  - Can cause transient forwarding loops



Convergence Delay

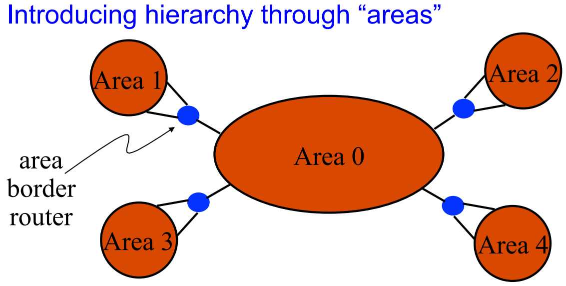
- Sources of convergence delay
  - Detection latency
  - Flooding of link-state information
  - Shortest-path computation
  - Creating the forwarding table
- Performance during convergence period
  - Lost packets due to blackholes and TTL expiry
  - Looping packets consuming resources
  - Out-of-order packets reaching the destination
- Very bad for VoIP, online gaming, and video

Reducing Convergence Delay

- Faster detection
  - Smaller hello timers
  - Link-layer technologies that can detect failures
- Faster flooding
  - Flooding immediately
  - Sending link-state packets with high-priority
- Faster computation
  - Faster processors on the routers
  - Incremental Dijkstra's algorithm
- Faster forwarding-table update
  - Data structures supporting incremental updates

Scaling Link-State Routing

- Overhead of link-state routing
  - Flooding link-state packets throughout the network
  - Running Dijkstra's shortest-path algorithm
- Introducing hierarchy through "areas"



Conclusions

- Routing is a distributed algorithm
  - React to changes in the topology
  - Compute the paths through the network
- Shortest-path link state routing
  - Flood link weights throughout the network
  - Compute shortest paths as a sum of link weights
  - Forward packets on next hop in the shortest path
- Convergence process
  - Changing from one topology to another
  - Transient periods of inconsistency across routers